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The Gaussian CI inference problem

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**MATHEMATISCHE
KOMPLEXITÄTSREDUKTION**

Gaussian conditional independence

Consider random variables $(\xi_i)_{i \in N} \sim \mathcal{N}(\mu, \Sigma)$. The *conditional independence (CI) statement* $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ conveys, informally, that if ξ_K is known, then learning the value of ξ_i does not give any information about ξ_j .

Definition

The polynomial $\Sigma[K] := \det \Sigma_{K,K}$ is a *principal minor* of Σ and $\Sigma[ij|K] := \det \Sigma_{iK,jK}$ is an *almost-principal minor*.

If Σ is positive-definite, then $\Sigma[K] > 0$, and $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ holds if and only if $\Sigma[ij|K] = 0$.



Almost-principal minors

$$\Sigma[ij] = x_{ij}$$

$$\Sigma[ij|k] = x_{ij}x_{kk} - x_{ik}x_{jk}$$

$$\Sigma[ij|kl] = x_{ij}x_{kk}x_{ll} - x_{il}x_{jl}x_{kk} + x_{il}x_{jk}x_{kl} + x_{ik}x_{jl}x_{kl} - x_{ij}x_{kl}^2 - x_{ik}x_{jk}x_{ll}$$

$$\begin{aligned}\Sigma[ij|klm] = & x_{ij}x_{kk}x_{ll}x_{mm} + x_{im}x_{jm}x_{kl}^2 - x_{im}x_{jl}x_{kl}x_{km} - x_{il}x_{jm}x_{kl}x_{km} + x_{il}x_{jl}x_{km}^2 \\& - x_{im}x_{jm}x_{kk}x_{ll} + x_{im}x_{jk}x_{km}x_{ll} + x_{ik}x_{jm}x_{km}x_{ll} - x_{ij}x_{km}^2x_{ll} \\& + x_{im}x_{jl}x_{kk}x_{lm} + x_{il}x_{jm}x_{kk}x_{lm} - x_{im}x_{jk}x_{kl}x_{lm} - x_{ik}x_{jm}x_{kl}x_{lm} \\& - x_{il}x_{jk}x_{km}x_{lm} - x_{ik}x_{jl}x_{km}x_{lm} + 2x_{ij}x_{kl}x_{km}x_{lm} + x_{ik}x_{jk}x_{lm}^2 \\& - x_{ij}x_{kk}x_{lm}^2 - x_{il}x_{jl}x_{kk}x_{mm} + x_{il}x_{jk}x_{kl}x_{mm} + x_{ik}x_{jl}x_{kl}x_{mm} \\& - x_{ij}x_{kl}^2x_{mm} - x_{ik}x_{jk}x_{ll}x_{mm}\end{aligned}$$

⋮



Gaussian CI models

Definition

A *CI constraint* is a CI statement $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ or its negation $\neg(\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K)$. They are *algebraic conditions* on the entries of Σ , equivalent to vanishing or non-vanishing of the almost-principal minors $\Sigma[ij|K]$.

Definition

The *model* of a set of CI constraints is the set of all positive-definite matrices which satisfy the constraints.



Gaussian CI models

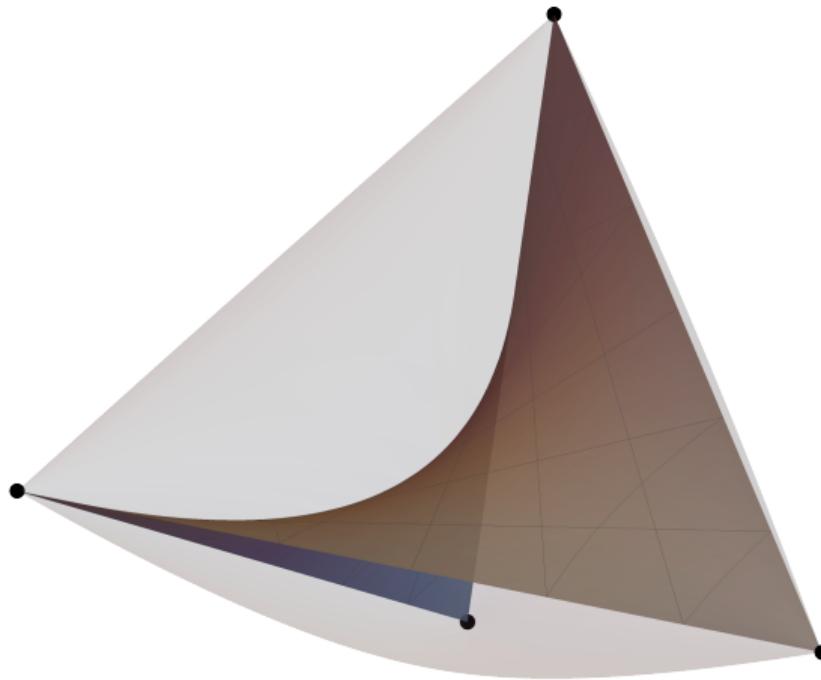


Figure: Gaussian model $\Sigma[12|3] = 0$ inside the ellotope.



Models and inference

Consider two sets of CI statements \mathcal{P} and \mathcal{Q} :

$$\bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q}$$



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$$\begin{array}{ccc} \bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q} & \iff & \mathcal{P} \cup \neg \mathcal{Q} \\ \text{is not valid} & & \text{has a model} \end{array}$$



Models and inference

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$$\begin{array}{ccc} \bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q} & \iff & \mathcal{P} \cup \neg \mathcal{Q} \\ \text{is not valid} & & \text{has a model} \end{array}$$

Reasoning about relevance statements in normally distributed random variables is **the same** as reasoning about the vanishing of very special kinds of determinants on very special kinds of varieties inside the positive-definite matrices.



Examples of CI inference

Consider a general positive-definite 3×3 correlation matrix

$$\Sigma = \begin{pmatrix} 1 & a & b \\ a & 1 & c \\ b & c & 1 \end{pmatrix}.$$

- If $\Sigma[12|3] = a - bc$ and $\Sigma[13|] = b$ vanish, then $\Sigma[12|] = a$ and $\Sigma[13|2] = b - ac$ must vanish as well:

$$\begin{aligned}(12|3) \wedge (13|) &\Rightarrow (12|), \\ (12|3) \wedge (13|) &\Rightarrow (13|2).\end{aligned}$$

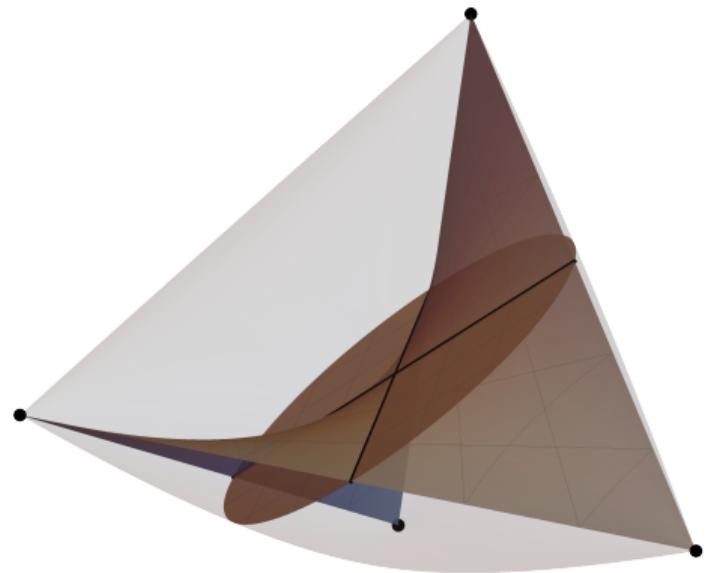


Examples of CI inference

$$\Sigma = \begin{pmatrix} 1 & a & b \\ a & 1 & c \\ b & c & 1 \end{pmatrix}$$

- If $\Sigma[12|] = a$ and $\Sigma[12|3] = a - bc$ vanish,
then $bc = \Sigma[13|] \cdot \Sigma[23|]$ must vanish:

$$(12|) \wedge (12|3) \Rightarrow (13|) \vee (23|).$$



No finite set of axioms

“There is no finite complete axiomatization of Gaussian CI”:

Theorem (Sullivant 2009)

As the matrix size n grows, there exist valid inference rules for Gaussians which need arbitrarily many antecedents.

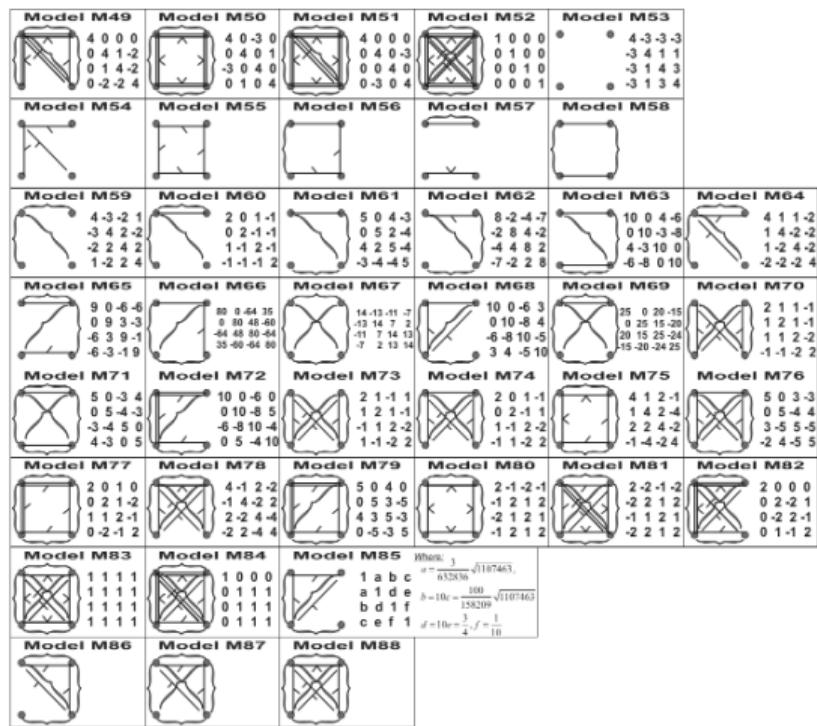
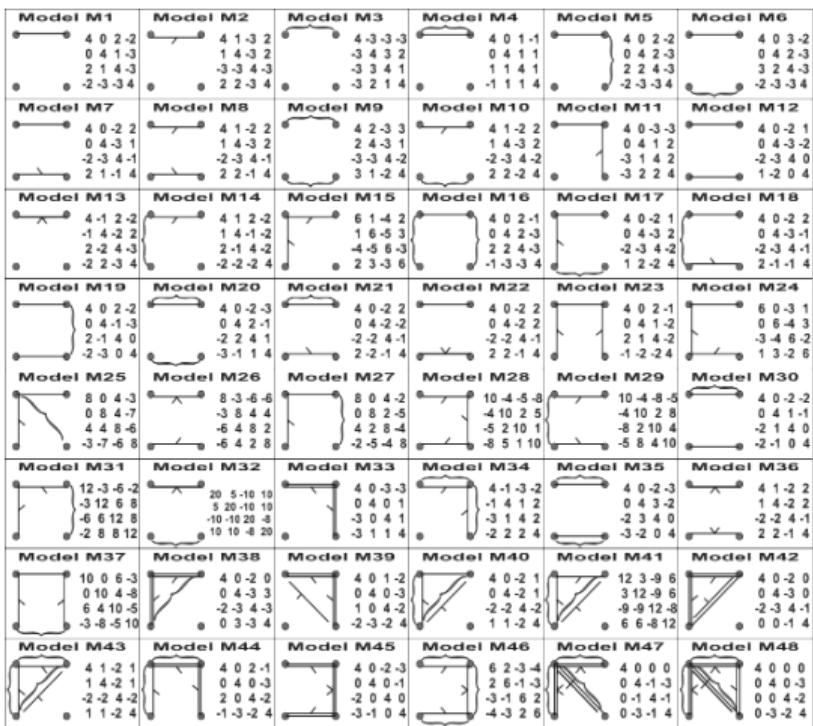
$$(12|3) \wedge (23|4) \wedge (34|1) \wedge (14|2) \Rightarrow (12|) \quad (n = 4)$$

$$(12|3) \wedge (23|4) \wedge (34|5) \wedge (45|1) \wedge (15|2) \Rightarrow (12|) \quad (n = 5)$$

$$(12|3) \wedge (23|4) \wedge (34|5) \wedge (45|6) \wedge (56|1) \wedge (16|2) \Rightarrow (12|) \quad (n = 6)$$

⋮





Šimeček's Question

Does every non-empty Gaussian CI model contain a rational point?

Model M85



Where: $a = \frac{3}{632836} \sqrt{1107463},$

$$b = 10c = \frac{100}{158209} \sqrt{1107463}$$

$$b = d = 10c$$

$$c = e = f = 1$$

$$d = 10e = \frac{3}{4}, f = \frac{1}{10}$$



Šimeček's Question

Does every non-empty Gaussian CI model contain a rational point?

Model M85



Where: $a = \frac{3}{632836} \sqrt{1107463},$

$$\begin{array}{cccccc} 1 & a & b & c \\ a & 1 & d & e \\ b & d & 1 & f \\ c & e & f & 1 \end{array}$$

$b = 10c = \frac{100}{158209} \sqrt{1107463}$

$$d = 10e = \frac{3}{4}, f = \frac{1}{10}$$

$$\begin{pmatrix} 357 & -21 & -343 & -147 \\ -21 & 357 & 119 & 51 \\ -343 & 119 & 357 & 153 \\ -147 & 51 & 153 & 357 \end{pmatrix}$$



Complexity bounds from real geometry

Theorem (Tarski's transfer principle)

If a polynomial system $\{f_i = 0, g_j > 0, h_k \neq 0\}$ has a solution over \mathbb{R} , then it has a solution in a finite real extension of \mathbb{Q} .

→ If $\bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q}$ is false, there exists a counterexample matrix Σ with algebraic entries.

Theorem (Positivstellensatz)

A polynomial F vanishes on the basic semialgebraic set $\{f_i = 0, g_j > 0, h_k \neq 0\}$ if and only if $0 \in \text{ideal}(f_i) + \text{cone}(g_j) + \text{monoid}^2(F, g_j, h_k)$.

→ If $\bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q}$ is true, there exists an algebraic proof for it with rational coefficients.



Universality theorems

Theorem (B. 2021)

*For every finite real extension \mathbb{K}/\mathbb{Q} there exists a Gaussian CI model $\mathcal{M}_{\mathbb{K}}$ such that:
for every \mathbb{L}/\mathbb{Q} , $\mathcal{M}_{\mathbb{K}}$ has an \mathbb{L} -rational point if and only if $\mathbb{K} \subseteq \mathbb{L}$.*

→ The answer to Šimeček's question is **NO**.

Theorem (B. 2021)

*The problem of deciding whether a CI inference formula is valid for all Gaussian distributions
is polynomial-time equivalent to the existential theory of the reals.*

→ Solving the inference problem can be used to check if arbitrary polynomial systems
have a solution over \mathbb{R} .



A bit of the proof idea

$\Sigma[ij] = x_{ij} \rightarrow$ impose $x_{kl} = x_{km} = x_{lm} = 0$ on a correlation matrix, then:

$$\begin{aligned}\Sigma[ij|klm] &= x_{ij}x_{kk}x_{ll}x_{mm} + x_{im}x_{jm}\underline{x_{kl}^2} - x_{im}x_{jl}\underline{x_{kl}x_{km}} - x_{il}x_{jm}\underline{x_{kl}x_{km}} + x_{il}x_{jl}\underline{x_{km}^2} \\ &\quad - x_{im}x_{jm}x_{kk}x_{ll} + x_{im}x_{jk}\underline{x_{km}x_{ll}} + x_{ik}x_{jm}\underline{x_{km}x_{ll}} - x_{ij}\underline{x_{km}^2}x_{ll} \\ &\quad + x_{im}x_{jl}x_{kk}\underline{x_{lm}} + x_{il}x_{jm}x_{kk}\underline{x_{lm}} - x_{im}x_{jk}\underline{x_{kl}x_{lm}} - x_{ik}x_{jm}\underline{x_{kl}x_{lm}} \\ &\quad - x_{il}x_{jk}\underline{x_{km}x_{lm}} - x_{ik}x_{jl}\underline{x_{km}x_{lm}} + 2x_{ij}\underline{x_{kl}x_{km}x_{lm}} + x_{ik}x_{jk}\underline{x_{lm}^2} \\ &\quad - x_{ij}x_{kk}\underline{x_{lm}^2} - x_{il}x_{jl}x_{kk}x_{mm} + x_{il}x_{jk}\underline{x_{kl}x_{mm}} + x_{ik}x_{jl}\underline{x_{kl}x_{mm}} \\ &\quad - x_{ij}\underline{x_{kl}^2}x_{mm} - x_{ik}x_{jk}x_{ll}x_{mm} \\ &= x_{ij} - \sum_{k=l,m} x_{ik}x_{jk} = x_{ij} - \left\langle \begin{pmatrix} x_{ik} \\ x_{il} \\ x_{im} \end{pmatrix}, \begin{pmatrix} x_{jk} \\ x_{jl} \\ x_{jm} \end{pmatrix} \right\rangle.\end{aligned}$$

The rest is 19th century projective geometry. Keyword: *von Staudt constructions*.



Approximations to the inference problem



Approximations to the inference problem

Theorem (Matúš 2005)

The following relations hold for every symmetric matrix Σ :

$$\Sigma[ij|L]^2 = \Sigma[iL] \cdot \Sigma[jL] - \Sigma[L] \cdot \Sigma[ijL]$$

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$



Approximations to the inference problem

Theorem (Matúš 2005)

The following relations hold for every symmetric matrix Σ :

$$\Sigma[ij|L]^2 = \Sigma[iL] \cdot \Sigma[jL] - \Sigma[L] \cdot \Sigma[ijL] \rightarrow \text{semimatroids}$$

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L] \rightarrow \text{gaussoids}$$

These relations define essential geometric properties of symmetric matrices in principal and almost-principal minor coordinates. Study their combinatorics!



The Gaussian CI configuration space

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

The *Gaussian CI configuration space* $\mathcal{G} \subseteq \mathbb{R}^{2^n} \times \mathbb{R}^{\binom{n}{2}2^{n-2}}$ consists of all vectors of principal and almost-principal minors of $\Sigma \in \text{PD}_n$.

$$\begin{pmatrix} 1 & a & b \\ a & 1 & c \\ b & c & 1 \end{pmatrix} \mapsto \begin{pmatrix} 1, 1, 1, 1, 1 - a^2, 1 - b^2, 1 - c^2, \\ 1 + 2abc - a^2 - b^2 - c^2, \\ a, a - bc, b, b - ac, c, c - ab \end{pmatrix} \in \mathbb{R}^{8+6}.$$



The Gaussian CI configuration space

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

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$$\begin{pmatrix} 1 & a & b \\ a & 1 & c \\ b & c & 1 \end{pmatrix} \mapsto \begin{pmatrix} 1, 1, 1, 1, 1 - a^2, 1 - b^2, 1 - c^2, \\ 1 + 2abc - a^2 - b^2 - c^2, \\ a, a - bc, b, b - ac, c, c - ab \end{pmatrix} \in \mathbb{R}^{8+6}.$$

Very wasteful encoding of a matrix, but this creates simple and useful relations on configuration vectors. The CI structure of Σ is encoded in the *zero pattern* of $c(\Sigma) \in \mathcal{G}$.



Combinatorial compatibility

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

Combinatorial compatibility means “fulfilling of relations under uncertainty”: What if we only knew that all $\Sigma[K] \neq 0$ and whether or not $\Sigma[ij|K] = 0$ for every $(ij|K)$?



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$$(ij|L) \wedge (ij|kL) \Rightarrow (ik|L) \vee (jk|L)$$

$$(ik|L) \wedge (ij|kL) \Rightarrow (ij|L)$$

⋮



Combinatorial compatibility

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

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$$(ij|L) \wedge (ij|kL) \Rightarrow (ik|L) \vee (jk|L)$$

$$(ik|L) \wedge (ij|kL) \Rightarrow (ij|L) \wedge (ik|jL)$$

$$(ij|kL) \wedge (ik|jL) \Rightarrow (ij|L) \wedge (ik|L)$$

$$(ij|L) \wedge (ik|L) \Rightarrow (ij|kL) \wedge (ik|jL)$$

This yields the definition of *gausroids*.



CI inference via SAT solvers

Since gaussoids have a finite axiomatization, a SAT solver like CaDiCaL can deduce implications under the gaussoid axioms:

$$\begin{aligned} & (12|3) \wedge (12|34) \wedge (24|1) \wedge (34|2) \\ \Rightarrow & (12|) \wedge (12|4) \wedge (24|) \wedge (24|3) \wedge (24|13) \wedge (34|) \end{aligned}$$

These conclusions are valid for all regular Gaussian distributions.



Oriented gausroids

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

What if we only knew that all $\text{sgn } \Sigma[K] = +1$ and the value of $\text{sgn } \Sigma[ij|K]$ for every $(ij|K)$?

$$+(ij|L) \wedge -(ij|kL) \Rightarrow [+(ik|L) \wedge +(jk|L)] \vee [-(ik|L) \wedge -(jk|L)]$$

→ *Oriented* and *orientable* gausroids.



Oriented gausoids

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

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$$+(ij|L) \wedge -(ij|kL) \Rightarrow [+(ik|L) \wedge +(jk|L)] \vee [-(ik|L) \wedge -(jk|L)]$$

→ *Oriented* and *orientable* gausoids.

$$(ij|L) \wedge (kl|L) \wedge (ik|jl) \wedge (jl|ik) \Rightarrow (ik|L)$$

$$(ij|L) \wedge (kl|il) \wedge (kl|jl) \wedge (ij|kl) \Rightarrow (kl|L)$$

$$(ij|L) \wedge (jl|kl) \wedge (kl|il) \wedge (ik|jl) \Rightarrow (ik|L)$$

$$(ij|kl) \wedge (ik|il) \wedge (il|jl) \Rightarrow (ij|L)$$

$$(ij|kl) \wedge (ik|il) \wedge (jl|il) \wedge (kl|jl) \Rightarrow (ij|L)$$



CI inference via SAT solvers II

Running the SAT solver CaDiCaL on the definition of oriented gausoids confirms that on their supports

$$(12|) \wedge (13|4) \wedge (14|5) \wedge (15|23) \wedge (23|5) \wedge (24|135) \wedge (34|12) \wedge (35|1) \wedge (45|2) \\ \Rightarrow \text{everything except } (25|K) \text{ for all } K.$$

Geometrically, the model is a Markov network.

Theorem (B. 2021+)

Orientable gausoids have no finite complete axiomatization.



Selfadhesivity

Gaussian multiinformation functions satisfy a *selfadhesivity* property just like entropy vectors of discrete random variables (Matúš 2007):

Theorem (B. 2020+)

Let Σ be a positive-definite matrix on \mathbb{R}^N and $K \subseteq N$. Let M be any set with $|M| = |N|$ and $M \cap N = K$. There exists a unique positive-definite Φ on $\mathbb{R}^{N \cup M}$ such that:

- $\Phi_N = \Sigma = \Phi_M$,
- $N \perp\!\!\!\perp M \mid K$ holds for Φ .



Structural selfadhesivity

Selfadhesivity can be formulated in CI terms and all necessary properties studied above can be **strengthened** by selfadhesivity. For example:



Structural selfadhesivity

Selfadhesivity can be formulated in CI terms and all necessary properties studied above can be **strengthened** by selfadhesivity. For example:

- \mathcal{O} is a *selfadhesive* orientable gaussoid on N if for every $K \subseteq N$ and M as above there exists an orientable gaussoid $\overline{\mathcal{O}}$ on $N \cup M$ such that
 - $\overline{\mathcal{O}}|_N = \mathcal{O} = \overline{\mathcal{O}}|_M$,
 - $(N, M|K) \in \overline{\mathcal{O}}$.

This property can be decided by 2^N calls to a program which decides orientability.

Every Gaussian CI structure is a selfadhesive orientable gaussoid, but not every orientable gaussoid is selfadhesive. Similarly for semimatroids.



The search for inference rules

Inference rules help characterize the *realizable* CI structures:

- 3-variate: 11 out of 64 by Matúš 2005.
- 4-variate: 629 out of 16 777 216 by Lněnička and Matúš 2007.
- 5-variate: *open!* (out of 1 208 925 819 614 629 174 706 176)
 - 254 826 gausroids modulo symmetry,
 - 87 792 of which are orientable semimatroids,
 - 84 434 of which are *selfadhesive* orientable semimatroids.



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Help wanted:

- Use linear programming and information inequalities.
- Non-linear information inequalities → Ahmadiyah and Vinzant 2021.
- Tropical approximations and valuated gausroids.
- Compute algebraic realization spaces.
- Find and certify real solutions to polynomial systems.



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